# Busting some database-myths episode 1



Erich Steiger Lead Database Architect



## Myths about me

- Lead Database Architect
  - 3.5 years at TWINT
- Software Engineering Background
  - Java / C++ and other programming languages
  - 27 years of Application development experience
  - Over 20 years in the financial sector
- Database knowledge
  - 27 years of Database experience
  - 21 years of Oracle experience
- First Program sold at age 15
- Unix since 1995 / Linux since 1993
- Private
  - 45 years old
  - Show-Jumping
  - Skiing



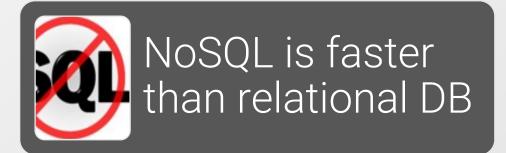
## **Myths about TWINT AG**

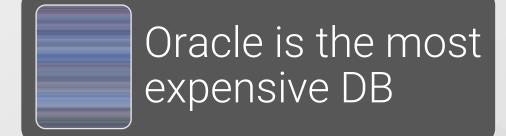
TWINT is Switzerland's market leading mobile payment provider

- Located in Zürich and Bern
- 210 Employees
- Around 70% of shops accept TWINT
- More than 5.5 Mio active users
- 38 Issuer Banks connected

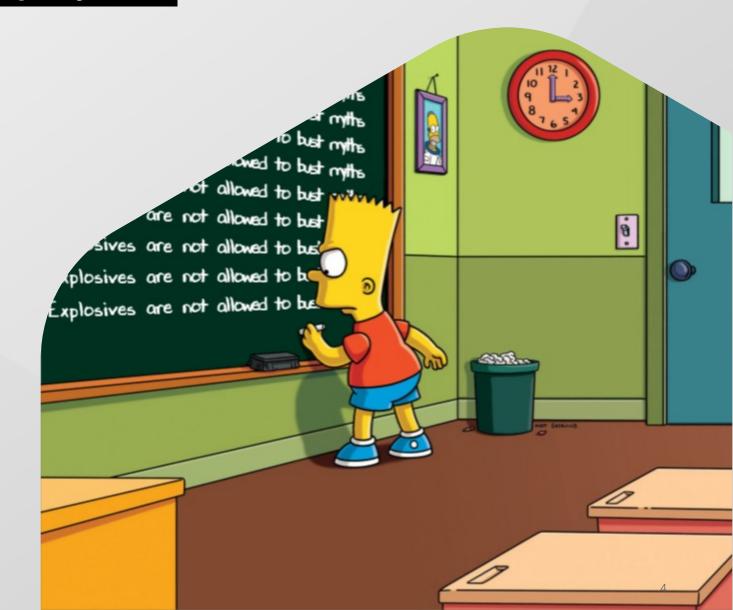


#### Who heard about the following Myths?



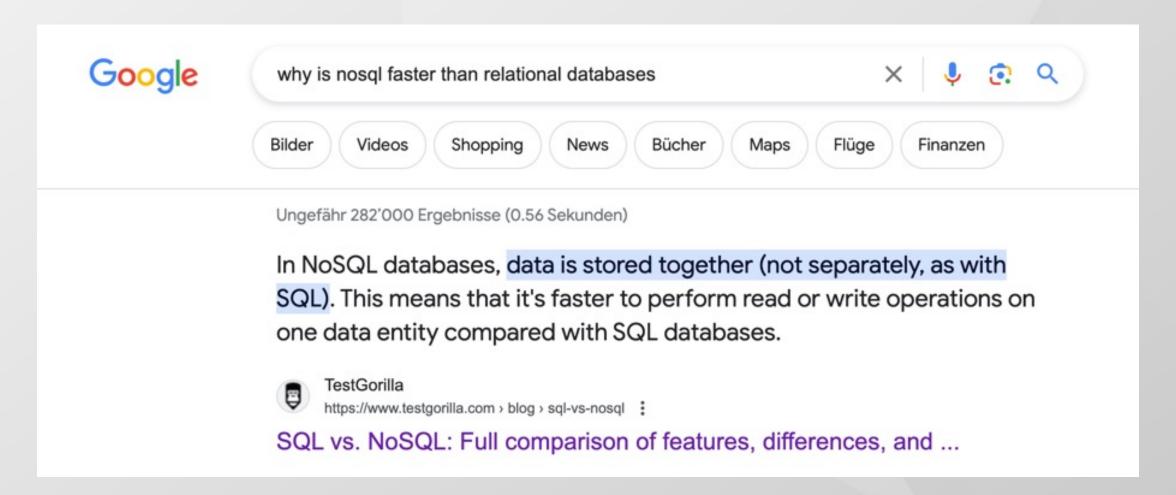




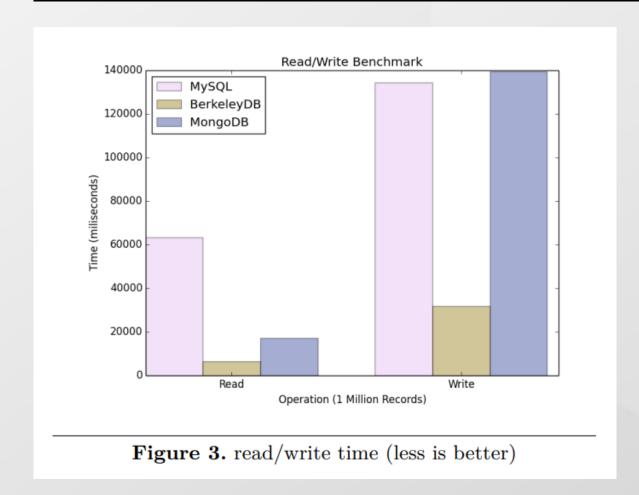


## NoSQL is faster than relational DB

## **Ask Google**



## Ruihan Wang, Zongyan Yang - University of Rochester





## Irrelevant in the Cloud



Hardware



Underlaying Technology Stack



Architecture

## Relevant in the Cloud



Performance and Latency



Throughput



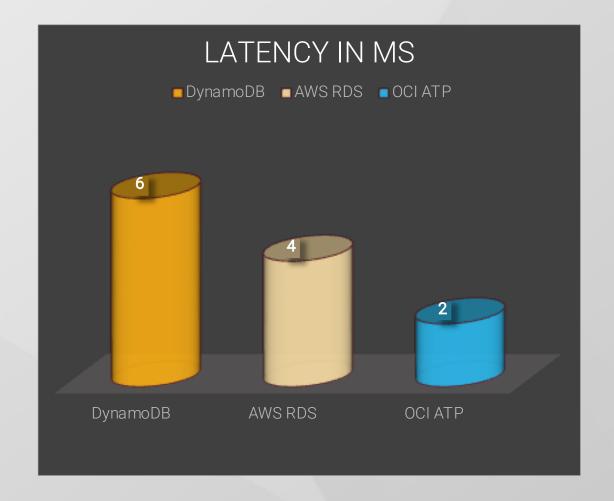
Price



Availability

#### RDBMS is faster than NoSQL

- Benchmark-Application inserts Business-Objects
- On Relational DBs a Business-Objects consists of 1 Main-Record and 2 Child-Records
- 22 indexes directly involved
- 5 FK constraints directly involved
- On DynamoDB one Business-Object is one record
- PK index
- SK index
- JSON Object instead of relational storage



#### Maybe lack of Knowledge?



Been using a relational database for a project recently. First time in a while that I've been using it for OLTP. Some quick reflections.

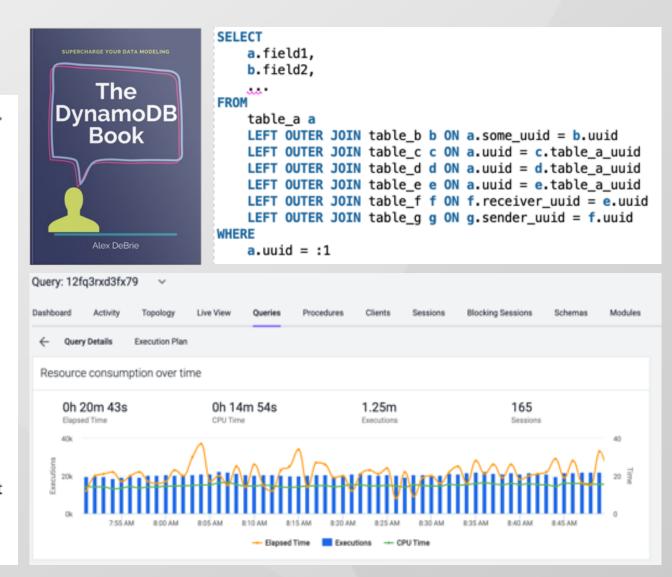
#### The good:

- The query flexibility! Ohh, the query flexibility. How I've missed thee. I really do appreciate that.

#### The less good:

- All the schema stuff is so annoying when iterating. Constantly going back + forth between doing it the right way and doing it the fast way.
- I get a tiny pit in my stomach not knowing what's happening on a given query and how it will work with more data. Especially true when you start having joins across 5 tables with some filter conditions.

2:26 PM · Mar 21, 2024 · 18.1K Views



# NoSQL is faster than relational DB

# Oracle is the most expensive DB

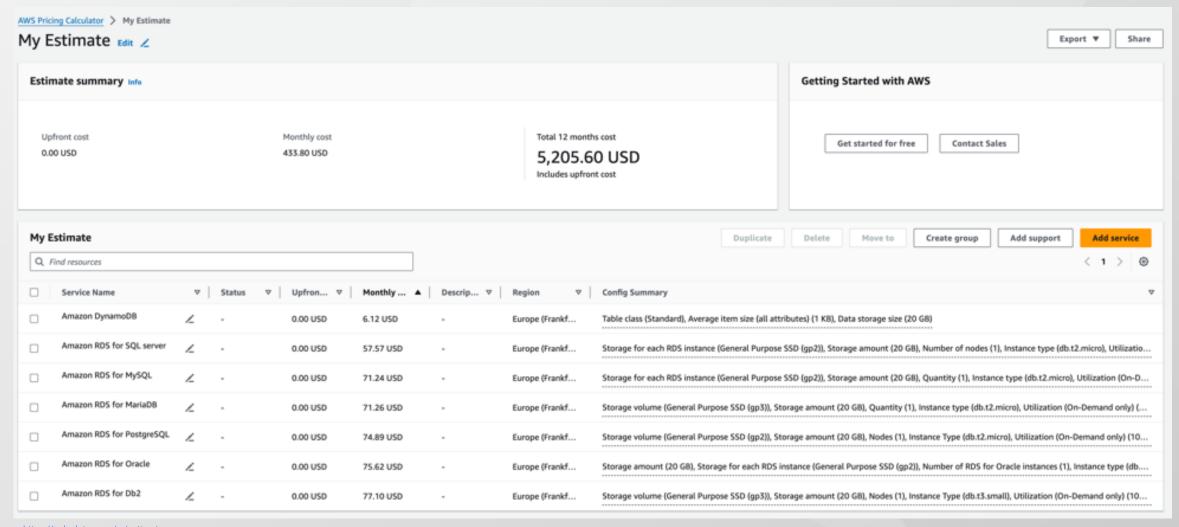
#### **DEMO**



**Warning:** the following figures are taken from cloud cost calculators between February and April 2024 in CHF. Please, be aware, that cloud prices and currency exchange rates may have changed since then. Also, the cost calculator software might have changed since then, sometimes on a daily base. The prices might not be 100% accurate anymore.

#### SOUG Day Spring, 17th April 2024

#### **AWS DB Cost estimate – smallest shapes**

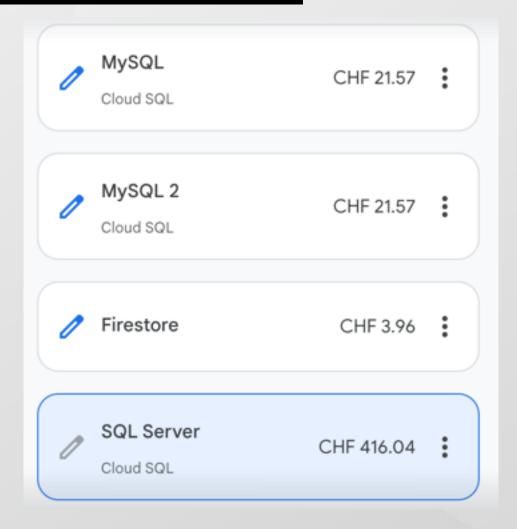


https://calculator.aws/ - /estimate

## Azure DB estimate – smallest shapes

| Azure Small Shapes                                |     |  |          |                   |                     |
|---|-----|--|----------|-------------------|---------------------|
|   |     |  |          |                   |                     |
| ✓ Azure SQL Database                              | (i) | Single Database, vCore, General Purpose, Provision   | <b>i</b> | Upfront: CHF 0.00 | Monthly: CHF 378.40 |
| <ul> <li>Azure Database for MySQL</li> </ul>      | i   | Single Server Deployment, Basic Tier, 1 Gen 5 (2 vC  | ÷ i      | Upfront: CHF 0.00 | Monthly: CHF 71.89  |
| <ul> <li>Azure Database for PostgreSQL</li> </ul> | i   | Single Server Deployment, Basic Tier, 1 Gen 5 (2 vC  | <b>i</b> | Upfront: CHF 0.00 | Monthly: CHF 71.89  |
| <ul> <li>Azure Database for MariaDB</li> </ul>    | i   | Basic Tier, 1 Gen 5 (2 vCore) x 730 Hours, 20 GB Sto | <b>+</b> | Upfront: CHF 0.00 | Monthly: CHF 71.89  |
| Azure Cosmos DB                                   | i   | Azure Cosmos DB for NoSQL (formerly Core), Stand     | <b>1</b> | Upfront: CHF 0.00 | Monthly: CHF 36.14  |
|   |     |  |          |                   |                     |

## **GCP DB estimate – medium shapes**

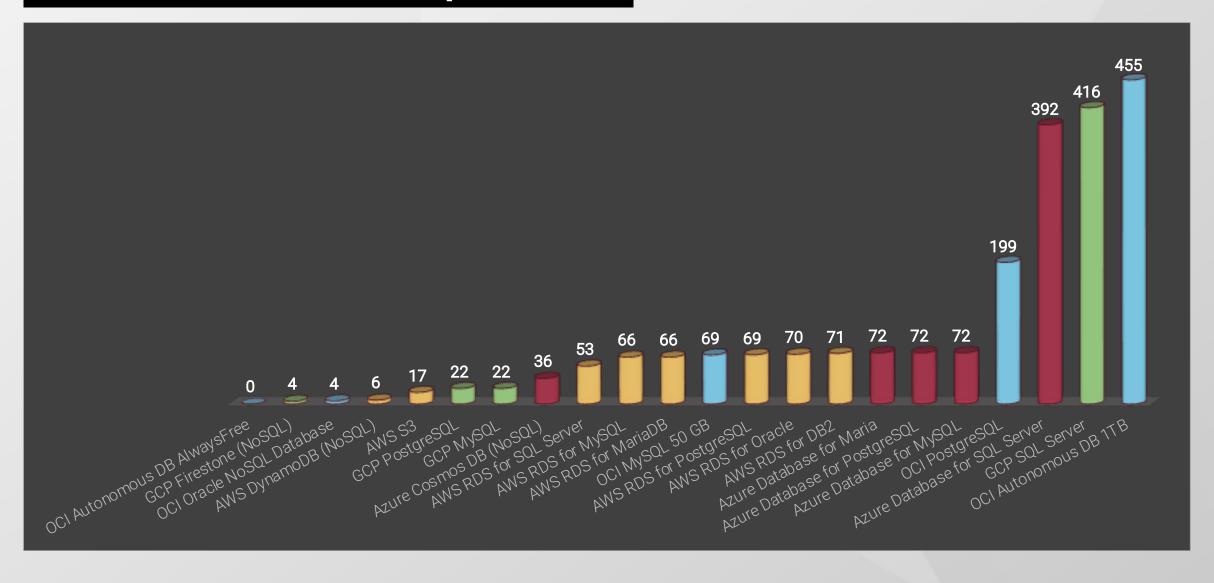


## OCI DB estimate – smallest shapes

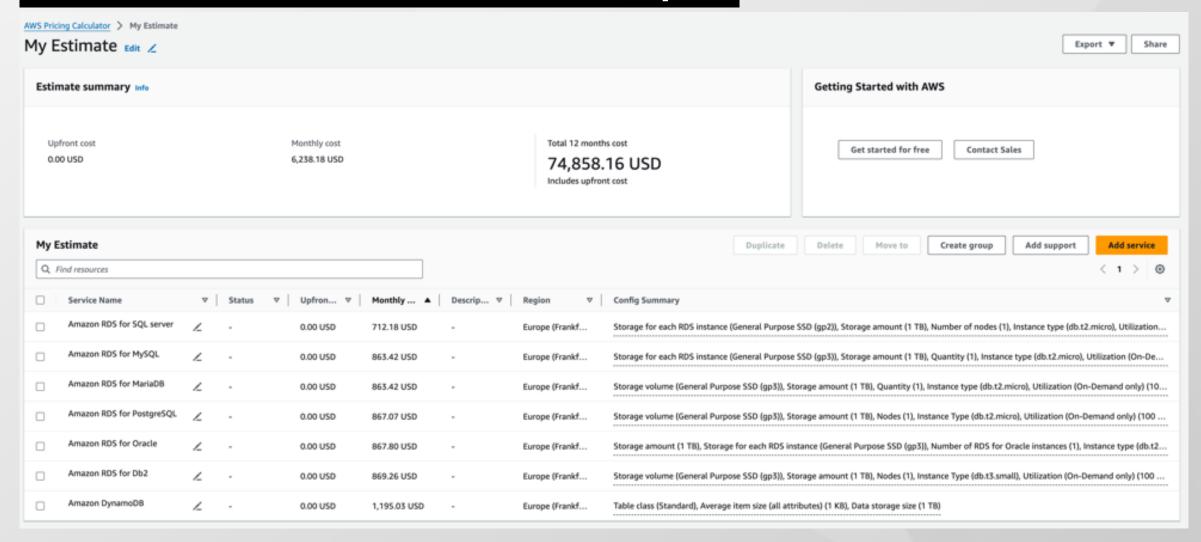
| Meine Schätzung                             | Gratis starten  CHF - Swiss Fran  Geschätzte Monatliche Kosten  CHF 729.13 | HH          |
|---|--|-------------|
| Konfiguration hinzufügen                    |  |             |
| Autonomous Database ···                     | Geschätzte Monatliche Kosten CHF 455.61                                    | <b> </b>    |
| MySQL Database Service ···                  | Geschätzte Monatliche Kosten CHF 69.35                                     | ; >         |
| Oracle NoSQL Database Cloud - On-Demand ··· | Geschätzte Monatliche Kosten CHF 4.49                                      | <b>&gt;</b> |
| Database with PostgreSQL ···                | Geschätzte Monatliche Kosten CHF 199.67                                    | 7 >         |

https://www.oracle.com/ch-de/cloud/costestimator.html16.4.2023

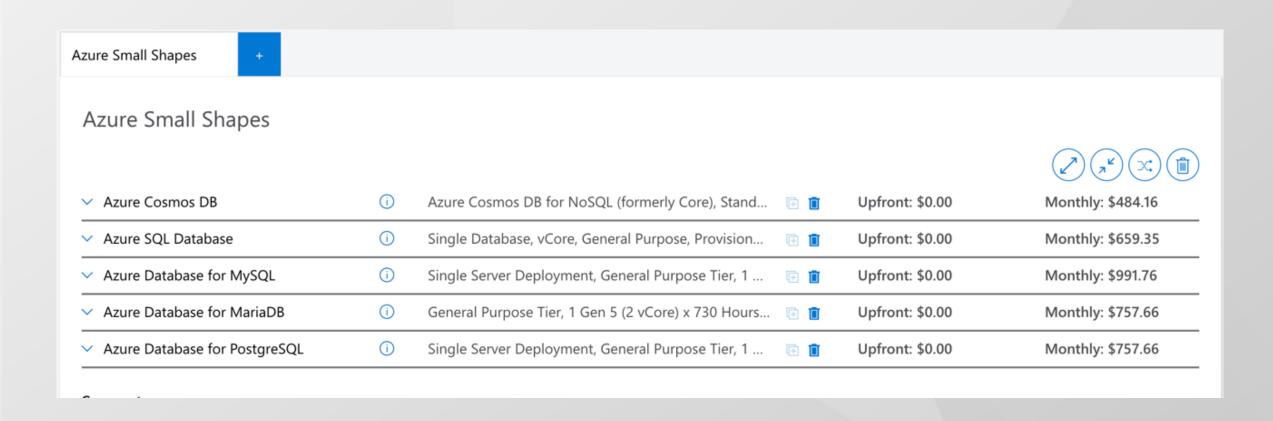
#### **Overview – smallest shapes 20 GB**



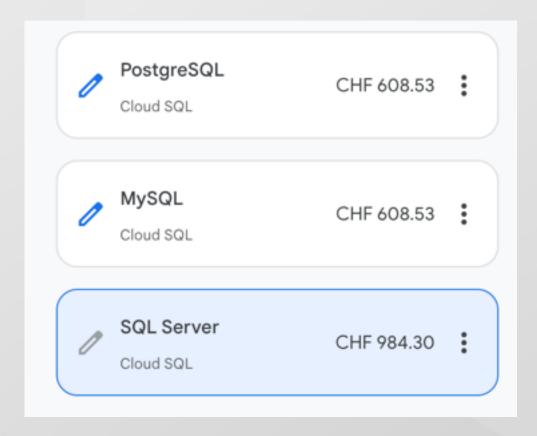
#### **AWS DB Cost estimate – medium shapes**



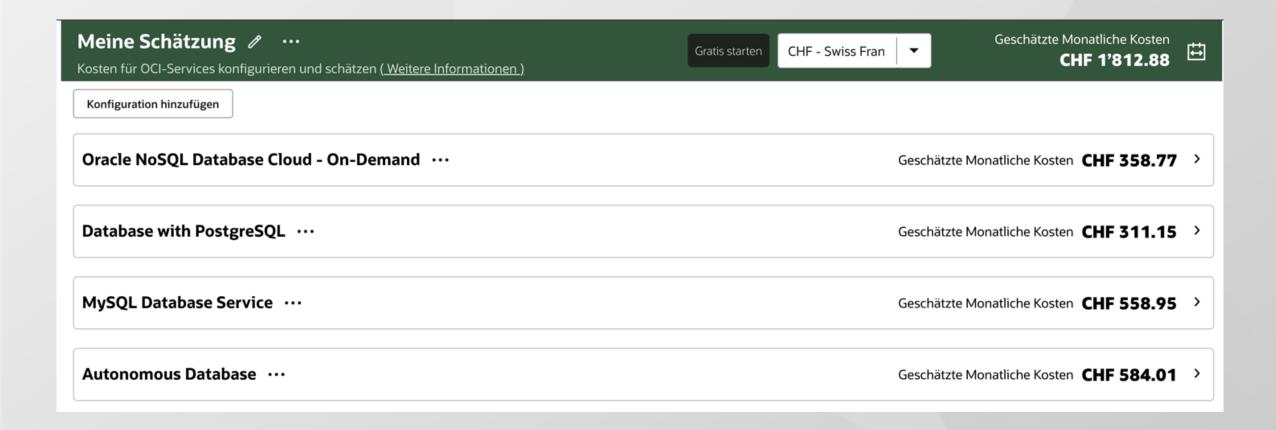
#### **Azure DB estimate – medium shapes**



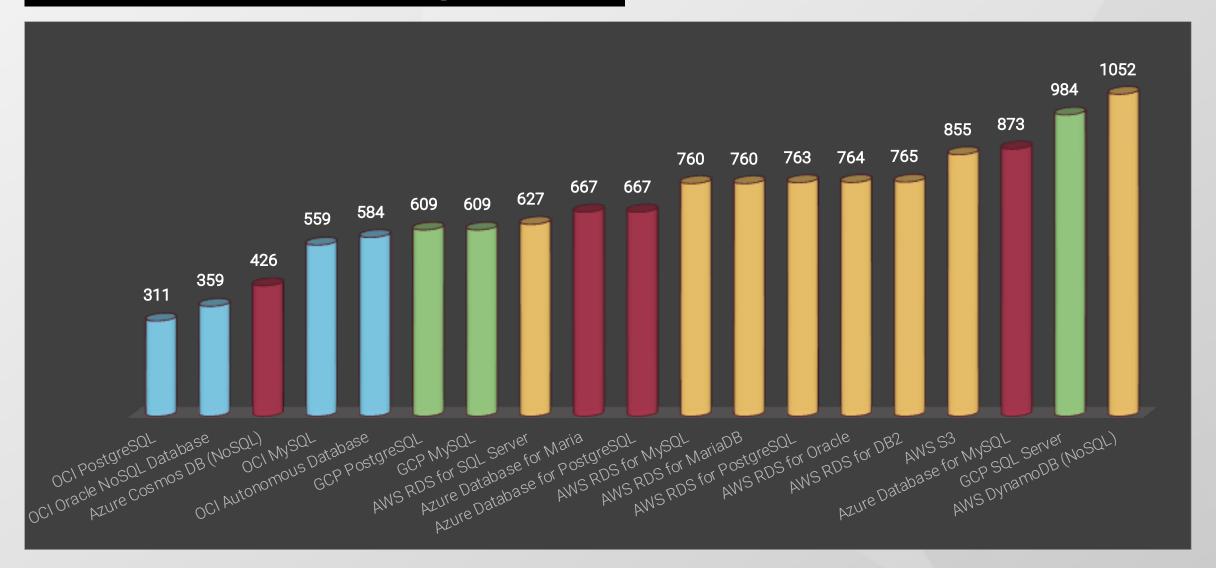
## **GCP DB estimate – medium shapes**



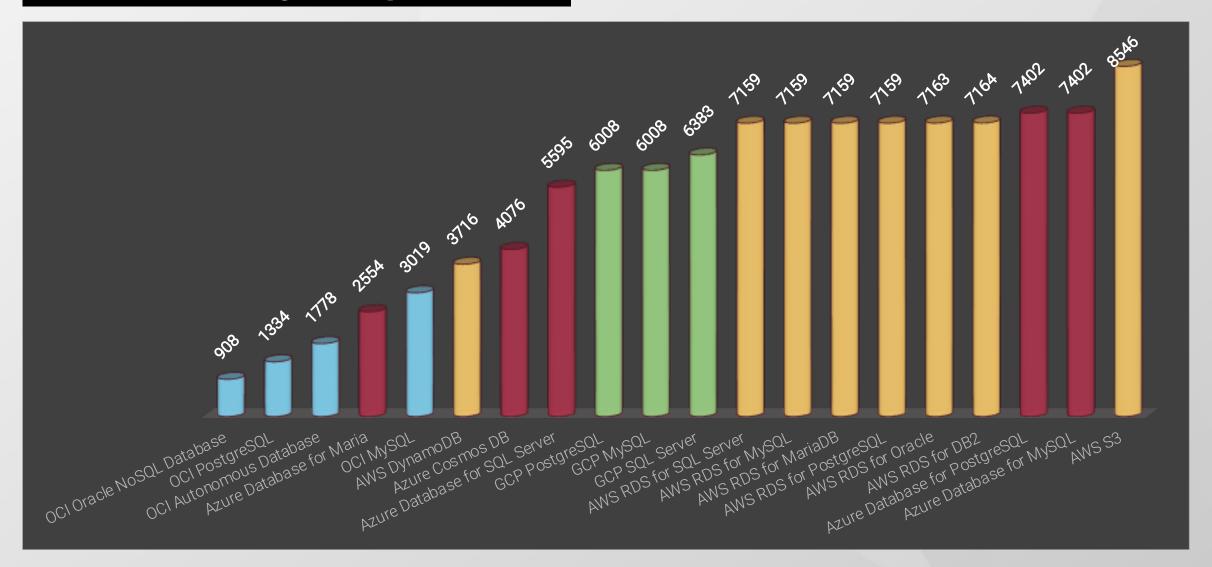
#### OCI DB estimate – medium shapes



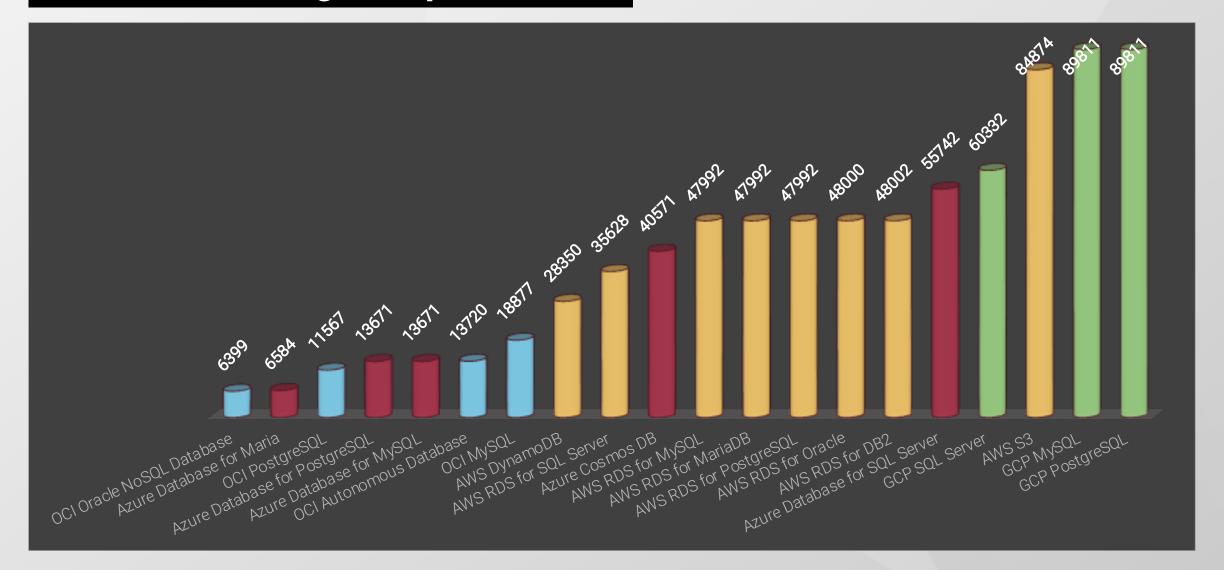
#### Overview – medium shapes 1 TB



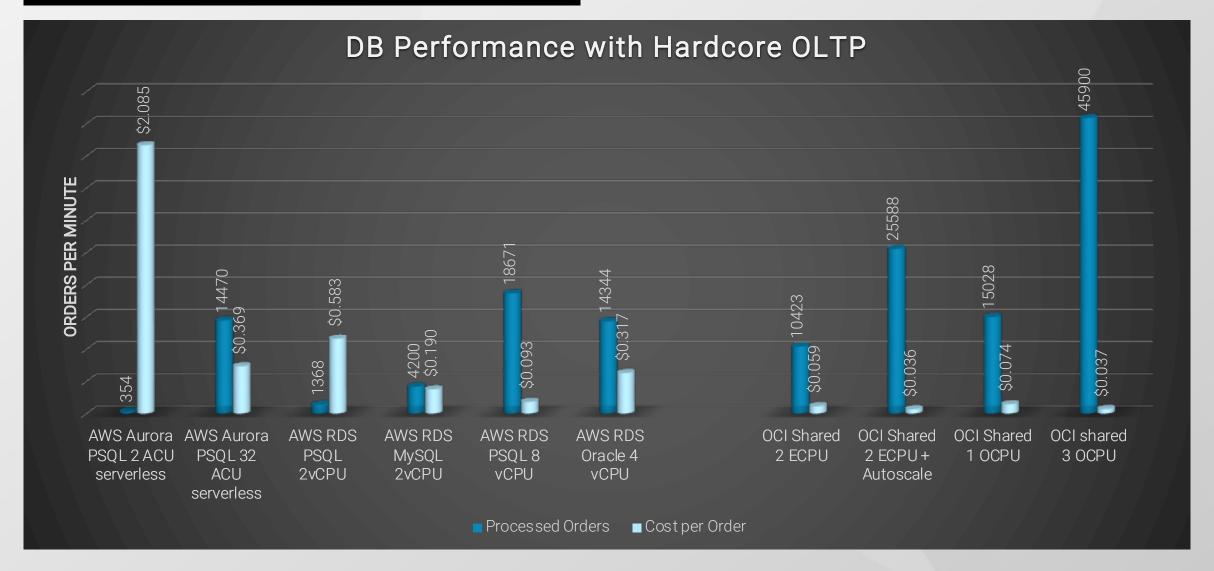
#### **Overview – large shapes 10 TB**



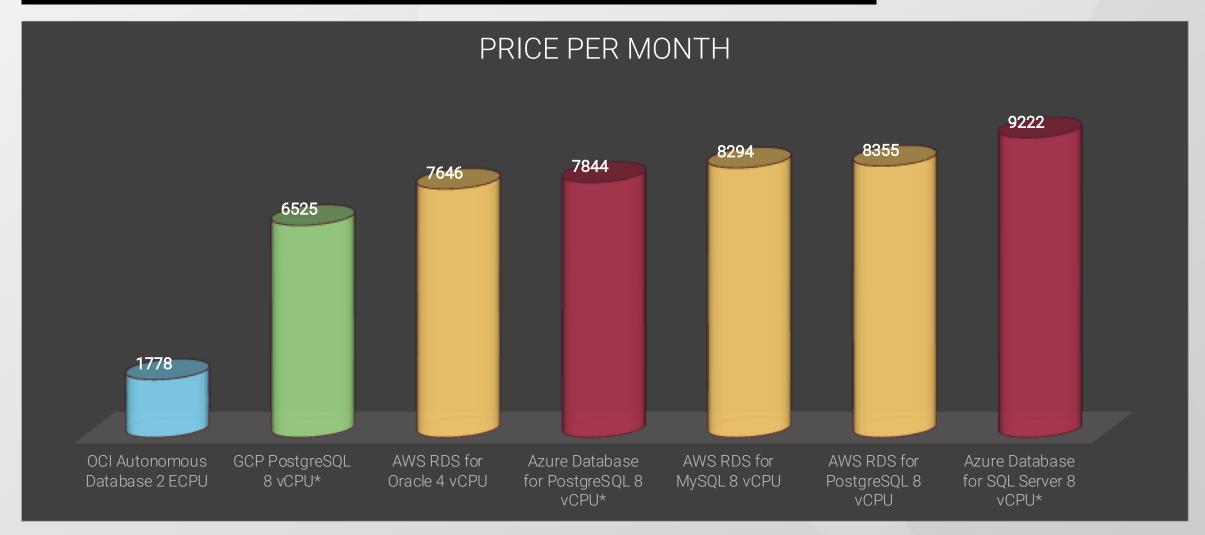
#### Overview - x-large shapes 100 TB



## **Throughput and price per Order**



## Databases able to process 10k orders per minute

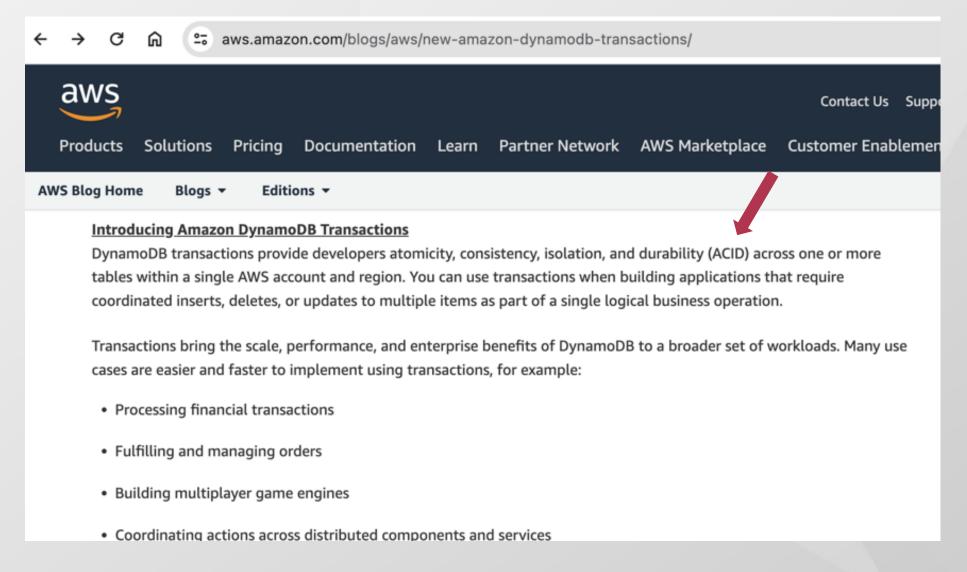


<sup>\* =</sup> no benchmark results taken, estimated size is guessed by results of AWS services and reality might look different

# Oracle is the most expensive DB

# **AWS DynamoDB is ACID-compliant**

## AWS says, DynamoDB is ACID compliant



## **DEMO 1 - DynamoDB**

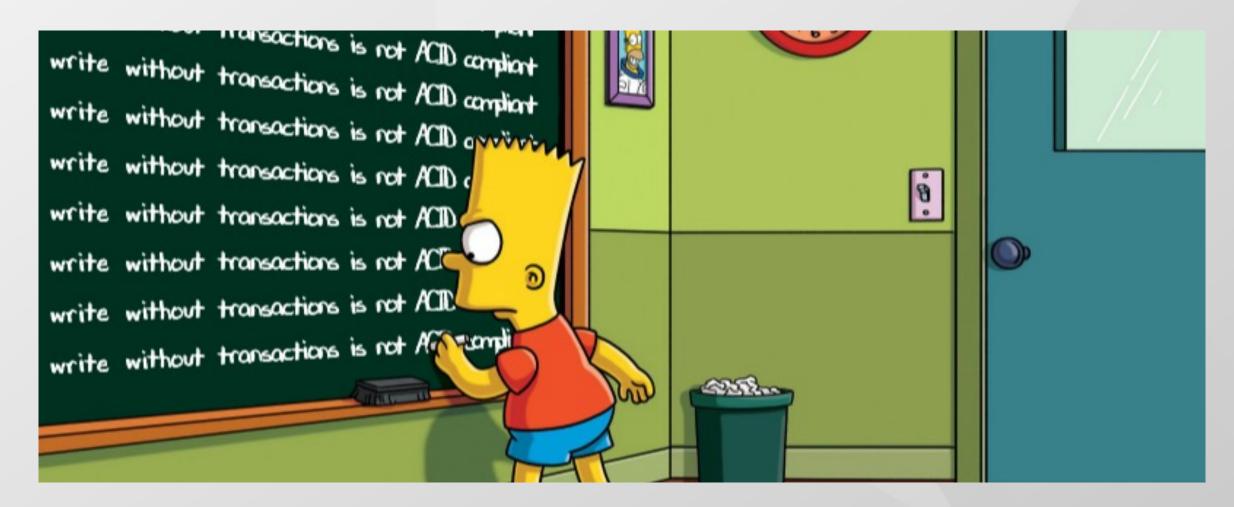
Write and read items

```
private void insertAndRead(Order order) {
    // save the item
    orderTable.putItem(order);
    LOGGER.debug("saved {}", order);

    // Retrieve the item
    Order orderRetrieved = orderTable.getItem(order);
    if (orderRetrieved == null) {
        insertStats.failedInserted += 1;
        LOGGER.error("Item retrieved: {} {}, stat: {}%", orderRetrieved, order, insertStats);
    } else {
        insertStats.successfulInserted += 1;
        LOGGER.debug("Item retrieved: {}", orderRetrieved);
    }
}
```



## Why did it fail?



#### DEMO 2 - DynamoDB

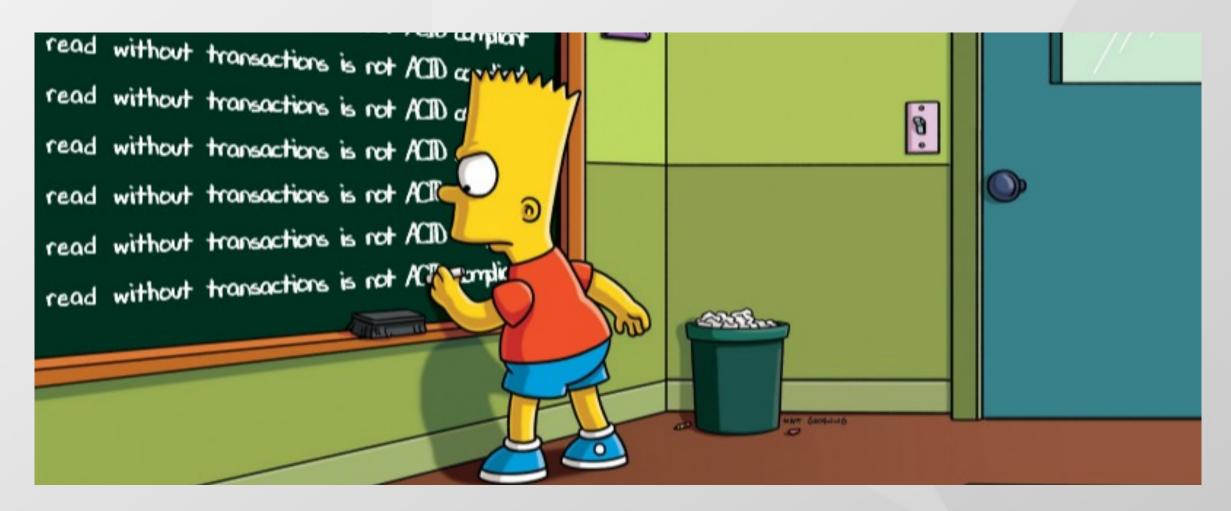
Write with transactions and then read the item

```
private void insertAndRead(Order order) {
    // save the item
    enhancedClient.transactWriteItems(b -> b.addPutItem(orderTable, order));
    LOGGER.debug("saved {}", order);

    // Retrieve the item
    Order orderRetrieved = orderTable.getItem(order);
    if (orderRetrieved == null) {
        insertStats.failedInserted += 1;
        LOGGER.error("Item retrieved: {} {}, stat: {}%", orderRetrieved, order, insertStats);
    } else {
        insertStats.successfulInserted += 1;
        LOGGER.debug("Item retrieved: {}", orderRetrieved);
    }
}
```



## Why did it fail?



## **DEMO 3 - DynamoDB**

Write with transactions and then read with transactions

## Why did it fail?



## DEMO 4 - DynamoDB

Write with transactions and then read with transactions and handle exceptions

```
private void insertAndRead(Order order) {
    // save the item
    enhancedClient.transactWriteItems(b -> b.addPutItem(orderTable, order));
    LOGGER.debug("saved {}", order);
    // Retrieve the item
        List<Document> orders = enhancedClient.transactGetItems(b -> b.addGetItem(orderTable, Key.builder()
                .partitionValue(order.getUuid().toString()).sortValue(order.getCreationTs().toString()).build()));
        if (orders.isEmpty()) {
            insertStats.failedInserted += 1;
            LOGGER.error("Item retrieved: {} {}, stat: {}%", orders.size(), order, insertStats);
            insertStats.successfulInserted += 1;
            LOGGER.debug("Item retrieved: {}", orders.get(0));
    } catch (TransactionConflictException e) {
        insertStats.failedInserted += 1;
        LOGGER.error("Exception while getting item: {}, stat: {}% {}", order, insertStats, e.getMessage());
        LOGGER.debug("Exception while get item: {}", order, e);
```



## Why did it fail?

Transactions might not be isolated

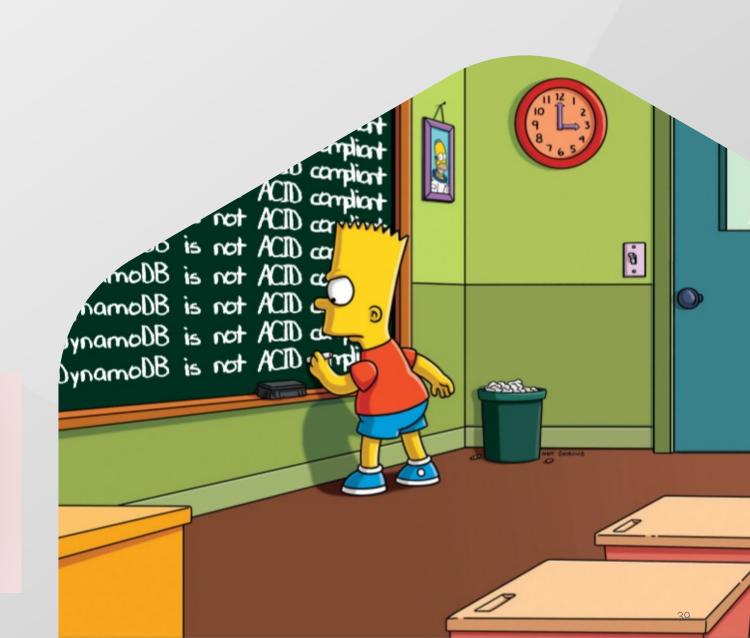
- I is broken

Read consistency not guaranteed

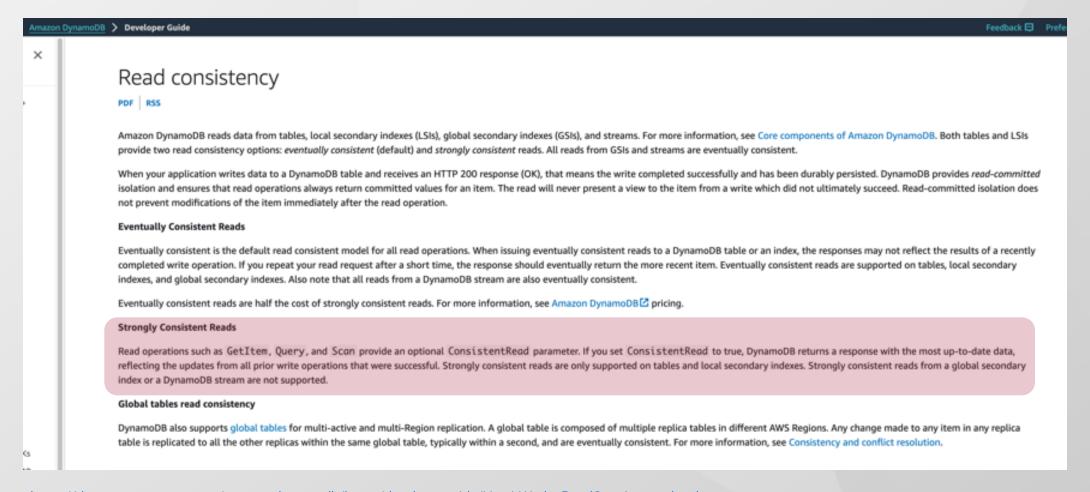
- C is broken

Eventual consistency and the way AWS implemented Transactions destroys ACID compliance

Read and Write cannot be combined within one transaction



#### AWS describes read-consistency as follows



https://docs.aws.amazon.com/amazondynamodb/latest/developerguide/HowItWorks.ReadConsistency.html

#### **Consistency Model description by AWS**

#### Consistency Model

PDF RSS

A database consistency model determines the manner and timing in which a successful write or update is reflected in a subsequent read operation of that same value.

Amazon DynamoDB lets you specify the desired consistency characteristics for each read request within an application. You can specify whether a read is eventually consistent or strongly consistent.

The eventual consistency option is the default in Amazon DynamoDB and maximizes the read throughput. However, an eventually consistent read might not always reflect the results of a recently completed write. Consistency across all copies of data is usually reached within a second.

A strongly consistent read in Amazon DynamoDB returns a result that reflects all writes that received a successful response prior to the read. To get a strongly consistent read result, you can specify optional parameters in a request. It takes more resources to process a strongly consistent read than an eventually consistent read. For more information about read consistency, see Data Read and Consistency Considerations.

Apache HBase reads and writes are strongly consistent. This means that all reads and writes to a single row in Apache HBase are atomic. Each concurrent reader and writer can make safe assumptions about the state of a row. Multi-versioning and time stamping in Apache HBase contribute to its strongly consistent model.

https://docs.aws.amazon.com/whitepapers/latest/comparing-dynamodb-and-hbase-for-nosql/consistency-model.html

#### Rochester Paper about ACID and CAP

#### 2.1. ACID Properties

We need to refer the ACID properties[12]:

#### Atomicity

A transaction is an atomic unit of processing; it should either be performed in its entirety or not performed at all.

#### Consistency preservation

A transaction should be consistency preserving, meaning that if it is completely executed from beginning to end without interference from other transactions, it should take the database from one consistent state to another.

#### Isolation

A transaction should appear as though it is being executed in iso-lation from other transactions, even though many transactions are execut- ing concurrently. That is, the execution of a transaction should not be interfered with by any other transactions executing concurrently.

#### Durability or permanency

The changes applied to the database by a com- mitted transaction must persist in the database. These changes must not be lost because of any failure.

#### **Rochester Paper about ACID and CAP**

#### 2.1. ACID Properties

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#### 2.2. CAP Theorem

For a distributed database, the CAP theorem states that it's impossible to simultaneously provide more than two out of the following three guarantees:

#### Consistency

Every read receives the most recent write or an error

#### Availability

Every request receives a (non-error) response – without guarantee that it contains the most recent write

#### Partition tolerance

The system continues to operate despite an arbitrary number of messages being dropped (or delayed) by the network between nodes

Based on CAP theorem, different database picks different combination of consistency, avalability, and partition tolerence:

- CA: Relational Database
- CP, AP: Non-Relational Database

# AWS Dynamed B is ACID-compliant

#### **Busted or confirmed?**



18/04/2024

Erich Steiger

# Busting some database-myths episode 1

Thanks for joining

# MYTH

**Erich Steiger** 

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#### Upcoming Talks

18<sup>th</sup> June 2024: DB Roundtable
Select AI chat what is TWINT?

